

Symbolic Reduction Technique for Real Power Flow Tracing

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Abstract— Power flow tracing is used to find the allocation of generators and loads in the line flows and generator contribution in loads and load contribution in generators. Due to multiplicity of solutions, recently optimal power flow tracing methods have been proposed. However, the size of optimization problem increases drastically with size of the system. Therefore, in this paper, we propose a symbolic reduction step wherein variables which will be set to zero at optimal are identified *a priori* by a graph theoretic analysis. This leads to reduce number of variables in the LP problem. Further, the reduced problem can be solved efficiently by sparse LP method. Results on large 488 node Indian utility system illustrate the necessity of proposed approach.

Keywords— Real power flow tracing, Depth first search, Variable reduction, Linear Programming (LP), Transmission system usage cost and loss allocation, Object oriented design

I. INTRODUCTION

Real power tracing is required to know the following components:

- How much of a particular load power has been supplied by each of generators in the system (Generation Allocation)?
- How much power is dispatched by a generator to a load (Load allocation)?
- What is the contribution of each generator and each load in the power flow of a line (Flow Allocation)?

All three contributions will be used to find the usage allocation. Usage allocation refers to the power contribution of each generator and/or each load. The advantage of knowing the usage allocation includes loss allocation associated to each path, cost assignment to transmission line pricing, congestion management, ancillary services and decision on scheduling of generators.

To compute the usage cost of a transmission, following methods have been proposed in literature [1].

- 1) *Postage stamp method*: In the postage-stamp method, transmission users are not differentiated by the “extent of use” of transmission facilities but charged based on an average embedded cost and the magnitude of transacted power.
- 2) *Contract path methodology*: Contract path method, on the other hand, assumes that the transacted power

would be confined to flow along an artificially specified path through the involved transmission systems. Once the contract path has been determined the transmission cost related to the specified path are assigned to the transaction. In reality, however, the actual path taken by a transaction may be quite different from the specified contract path thus involving the use of transmission facilities outside the contracted systems.

- 3) *MW-Mile method*: MW-Mile methodology may be regarded as the first pricing strategy proposed for the recovery of fixed transmission costs based on the actual use of transmission network. In this method, charges for each wheeling transaction are based on the measure of transmission capacity utilized in this transaction. However, it is difficult to segregate flow in a line in components of various simultaneous transactions.
- 4) *Proportionate sharing*: This method was proposed by J. Bialek [2]. The procedure assumes that the power reaching a certain node in the electric network is proportionately shared by all the paths going out from that node.

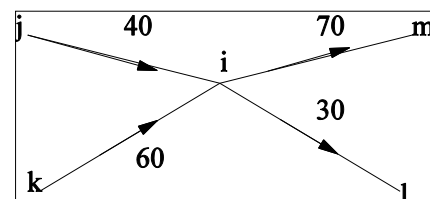


Fig. 1. Proportional sharing principle.

The principle of “proportional sharing” illustrated in Fig. 1 where four lines are connected to node *i*, two with inflows and two with outflows. The total flow through the node is $40 + 60 = 100$ units of which 40% is supplied by line *j – i*, and 60% by line *k – i*. As electricity is indistinguishable and each of the outflows down the line from node *i* depends only on the voltage gradient and the line impedance, it was assumed that each unit leaving the node contains the same proportion of the inflows as the total nodal flow. Hence the 70 units outflowing in line *i – m* consists of $\frac{40 \times 70}{100} = 28$ supplied by line *j – i* and $\frac{60 \times 70}{100} = 42$ supplied by line *k – i*.

Similarly, the 30 units outflowing in line *i – l* consists of $\frac{40 \times 30}{100} = 12$ supplied by line *j – i* and $\frac{60 \times 30}{100} = 18$ supplied by

line $k - i$. Even though the *proportional sharing* assumption cannot be proved or disproved, its application leads to simplicity.

- 5) Tracing compliant modified postage stamp method: Since, tracing problem is amenable to multiple solutions; we need to define an objective to sample out the appropriate solution. Therefore tracing problem has been framed as an optimization problem in [3]. The objective here is to choose the tracing compliant solution that is nearest to the postage stamp allocation of transmission fixed costs. Thus, advantages of usage based and nearly equitable distribution of transmission fixed cost allocation is achieved simultaneously. Same methodology is used for the loss allocation also.
- 6) Min-max fairness method [4], [5]: A Min-max fair power flow tracing has been proposed in [4]. The min-max power flow tracing further improved in [5] for application in large systems. A solution is said to be min-max fair, if any reduction in, say, percentage usage costs (MU/MW) for one entity leads to increase in share of another entity, which has either equal or higher percentage cost. In absence of traceability constraint min-max fairness and postage stamp method would lead to identical cost distribution. An advantage of min-max fairness model over modified postage stamp method is that it guarantee's fairness for each individual user, as within traceable constrained space, no other solution exists by which a higher cost entity (MU/MW) can off load usage costs to a lower cost entity. Min-max fairness also meets an important fairness criterion known as 'aggregate invariance' which implies that artificial splitting of load or generator at a bus will not alter the final solution.

While optimal tracing is superior to normal tracing as one can rigorously impose a fairness related objective on the space of traceable solutions, there are certain computational challenges to be overcome. In particular number of variables in optimal tracing problem is increased drastically. Thus, as the system size increases, the optimization solver becomes inefficient.

In this paper, we proposed a variable reduction technique. The rest of the paper is organized as follows: the optimal tracing methods i.e. Tracing compliant modified postage stamp method and Min-max fairness tracing method are explained in Section II. Symbolic reduction technique to improve computational efficiency of optimal tracing problem is explained in Section-III. Object oriented design in unified modelling language (UML) of tracing engine is discussed in Section IV. Section V presents the results and Section VI concludes the paper.

II. OPTIMIZATION APPROACH TO REAL POWER TRACING: AN APPLICATIONS TO TRANSMISSION FIXED COST ALLOCATION [3]

The optimal tracing problem can be compactly defined as follows:

$$\text{Problem OPT}(x, y): \min_{\{x, y\} \in S} f(x, y) \quad (1)$$

The set S represents the set of all possible tracing solutions and a specific set of x and y vectors represents a solution to generation and load tracing problem. It is shown that set can be characterized by a set of linear equality and inequality

constraints. In fact, set is both compact and convex. This leads to a linear constrained optimization problem. It models relationship between the flow entities and associated network usage costs. Constraints in **OPT** problem can be grouped into the following specifications

- flow specification constraints for series branches, i.e., transmission lines and transformers;
- source and sink specification constraints pertaining to shunts, e.g., generators and loads;
- conservation of commodity flow constraints.

A. Flow Specification Constraints

Traditionally, two types of tracing problems, viz., generation tracing and load tracing, are discussed in the literature. Generation tracing traces generator flows to loads, while load tracing traces load flows to generators. We first discuss modeling of the flow specification constraints for generation tracing problem.

1. *Generation Tracing*: Let P_{lm} (MW) be the flow on a line lm . Flow is supplied by generators G_1, G_2, \dots, G_{ng} with components $P_{lm}^{G_1}, P_{lm}^{G_2} \dots P_{lm}^{G_{ng}}$

$$P_{lm} = P_{lm}^{G_1} + P_{lm}^{G_2} + \dots + P_{lm}^{G_{ng}} \quad (2)$$

The component of generator G_k on lm line can be expressed as fraction x_{lm}^k of the total injection by generator G_k , i.e., . Therefore

$$P_{lm}^{G_k} = x_{lm}^k \cdot P_{G_k} \quad (3)$$

$$P_{lm} = \sum_{k=1}^{ng} x_{lm}^k \cdot P_{G_k}, \forall \text{ set of lines} \quad (4)$$

Since the branch flows are known and x are unknown, flow equations for generation allocation can be written as follow:

$$[A_{flowd}][x_{flow}] = [b_{flowd}] \quad (5)$$

2. *Load Tracing*: The power flow P_{lm} on line lm can also be expressed as a summation of load components, i.e.,

$$P_{lm} = P_{lm}^{L_1} + P_{lm}^{L_2} + \dots + P_{lm}^{L_{nL}} \quad (6)$$

The component of load (P_{Li}) on line lm is expressed as a fraction y_{lm}^i of load P_{Li} as follows:

$$P_{lm}^{L_i} = y_{lm}^i \cdot P_{L_i} \quad (7)$$

$$P_{lm} = \sum_{i=1}^{nL} y_{lm}^i \cdot P_{L_i}, \forall \text{ set of lines} \quad (8)$$

In matrix form, the flow equations for load allocation can be written as follows:

$$[A_{flowu}][y_{flow}] = [b_{flowu}] \quad (9)$$

B. Source and Sink Specification Constraints

1. *Generation Tracing*: In a generation tracing problem, it is necessary to write sink (load) constraints. They express the contribution of generators in loads. For a

load (P_{L_i}), the contribution of various generators is governed by the following constraint:

$$P_{L_i} = P_{L_i}^{G_1} + P_{L_i}^{G_2} + \dots + P_{L_i}^{G_{n_G}} \quad (10)$$

$$P_{L_i} = \sum_{k=1}^{n_G} x_i^k \cdot P_{G_k}, i = 1, \dots, n_L \quad (11)$$

In the matrix form, the load equations for generation allocation can be written as follows:

$$[A_{inj_d}][x_{inj}] = [b_{inj_d}] \quad (12)$$

2. *Load Tracing*: In the load tracing problem, it is necessary to model the share of loads in a generator. Let

$$P_{G_k} = P_{G_k}^{L_1} + P_{G_k}^{L_2} + \dots + P_{G_k}^{L_{n_L}} \text{ i.e., } P_{G_k} = \sum_{i=1}^{n_L} y_k^i P_{L_i}. \quad (13)$$

In matrix form, the generator equations for load allocation can be written as follows:

$$[A_{inj_u}][y_{inj}] = [b_{inj_u}] \quad (14)$$

C. Conservation of Commodity flow constraints

The conservation of flow constraints can be neatly expressed by using arc or bus incidence matrix M of the underlying graph. In the matrix, rows correspond to nodes and columns to arcs. The entry $M(i, j)$ is set to 1 if arc is outgoing at node; it is -1 if the arc is incoming at node; else, it is set to zero. The shunt arcs have one node as ground, which is not modeled in M . The corresponding entry in M is either 1 or -1, depending upon whether the arc represents load or generation.

- 1) *Generation Tracing*: Let M_d represents sub-matrix formed by considering series branches and shunt loads

$$M_d = [M_{nb}, M_L] \quad (15)$$

$$[M_d][x^k] = [e_k] \quad k = 1 \dots n_G \quad (16)$$

Where x^k represents the set of x -variables for lines and loads associated with the k^{th} generator. \bar{k} is the node at which k^{th} generator is connected, and e_k is the k^{th} column of identity matrix.

Instead of partitioning x variables by generator numbers, they can be partitioned by series branch flow (x_{flow}) and shunt flow x_{inj} variables. The set of the continuity equations can be rearranged and written in block matrix notations with and the variable partition as follows:

$$[A_{cont_{flow_d}} \quad A_{cont_{inj_d}}] \begin{bmatrix} x_{flow} \\ x_{inj} \end{bmatrix} = [b_{cont_d}] \quad (17)$$

- 2) *Load Tracing*: The following similar arguments as in generation tracing, the continuity equations for load tracing are given as follows:

$$[M_u][y^i] = [e_i] \quad i = 1 \dots n_L \quad (18)$$

where \bar{i} is the node at which the i^{th} load is connected.

$$[A_{cont_{flow_u}} \quad A_{cont_{inj_u}}] \begin{bmatrix} y_{flow} \\ y_{inj} \end{bmatrix} = [b_{cont_u}] \quad (19)$$

D. Lossy Flow network

For a lossy flow network, the flow reduces along the arc from the sending end to the receiving end. Under such situations, it is necessary to model two flow equations per line, one for the sending end and one for the receiving end. An alternative approach has been developed to restrict the number of variables and equations in the optimization problem to the corresponding lossless formulation. For taking care of lossy flow network we need to change the M_d and M_u . Let o_i is origin and d_i is destination of an arc i .

$$M_d^{loss}(o_i, i) = 1 \quad (20)$$

$$M_d^{loss}(d_i, i) = -\frac{P_{lm}^r}{P_{lm}^s}, \forall lm \in \text{set of lines} \quad (21)$$

$$M_u^{loss}(o_i, i) = \frac{P_{lm}^s}{P_{lm}^r}, \forall lm \in \text{set of lines} \quad (22)$$

$$M_u^{loss}(d_i, i) = -1 \quad (23)$$

E. Modeling of Boundary Conditions

In a megawatt flow network, if the results are consistent, then this difference, i.e., power dispatched by generator k to load i minus power received from generator k by load i , should correspond to the loss incurred in the generator- k load- i interaction. Let us define this loss as $loss_i^k$, where

$$loss_i^k = y_k^i P_{L_i} - x_i^k P_{G_k} \quad (24)$$

The above equation should satisfy the following loss properties of a network:

$$y_k^i P_{L_i} - x_i^k P_{G_k} \geq 0, \forall k \in \{1, 2, \dots, n_G\} \quad (25)$$

$$P_{loss} = \sum_{k=1}^{n_G} \sum_{i=1}^{n_L} y_k^i P_{L_i} - x_i^k P_{G_k} \quad (26)$$

Equation (25) models the constraint that loss is a non-negative number, while (26) models the requirement that total generator-load interaction losses should be equal to the power loss in the system.

F. Explicit formulation

1) *Objective Function for Transmission Fixed Cost Allocation Using Modified Postage Stamp Method*: The transmission system usage cost per MW paid by a load can be worked out as follows:

$$trprice_{Load_{pu}}^i = \frac{\sum_{lm} y_{lm}^i P_{L_i} \bar{c}_{lm} \alpha}{P_{L_i}} = \sum_{lm} y_{lm}^i \bar{c}_{lm} \alpha \quad (27)$$

The transmission system usage cost per MW paid by a generator can be worked out as follows:

$$\begin{aligned} trprice_{Gen_{pu}}^k &= \frac{\sum_{\forall lm} x_{lm}^k P_{G_k} \bar{c}_{lm} (1 - \alpha)}{P_{G_k}} \\ &= \sum_{\forall lm} x_{lm}^k \bar{c}_{lm} (1 - \alpha) \end{aligned} \quad (28)$$

where \bar{c}_{lm} is the cost of line per MW, α is the P. U. share of the load in total transmission cost, $0 \leq \alpha \leq 1$. Further, we assumed without loss of generality that $TSC_{net} = \sum_{\forall lines} \bar{c}_{lm} P_{lm}$.

Let the postage stamp rate for the loads be given by $k_{Load_T}^*$. Then

$$k_{Load_T}^* = \frac{\sum_{for all lm} P_{lm} \bar{c}_{lm} \alpha}{\sum_{i=1}^{n_L} P_{L_i}} \quad (29)$$

Let the postage stamp rate for the loads be given by $k_{Gen_T}^*$. Then

$$k_{Gen_T}^* = \frac{\sum_{\forall lm} P_{lm} \bar{c}_{lm} (1 - \alpha)}{\sum_{k=1}^{n_G} P_{G_k}} \quad (30)$$

The aim of tracing compliant postage stamp method is to compute the closest traceable solution to the proportionate distribution of transmission system usage costs. Therefore, the objective function $f\{x, y\}$ in (1) is written as

$$\begin{aligned} f\{x, y\} &= \sum_{i=1}^{n_L} |trprice_{Load_{pu}}^i - k_{Load_T}^*| \\ &+ \sum_{k=1}^{n_G} |trprice_{Gen_{pu}}^k - k_{Gen_T}^*| \end{aligned} \quad (31)$$

2) *Objective Function for Loss Allocation:* Another choice of objective function can be developed on similar lines for equitable distribution of losses. The per unit loss for load is given as follows:

$$loss_{pu}^i = \left(\sum_{k=1}^{n_G} y_k^i \right) - 1 \quad (32)$$

The per unit loss for generator is given as follows:

$$loss_{pu}^k = 1 - \left(\sum_{i=1}^{n_L} x_i^k \right) \quad (33)$$

Let us assume that percentage share of loss for generators in total loss of system is ShrLossG and percentage share of loss for loads in total loss of system is ShrLossL.

Therefore the objective function $f\{x, y\}$ in (1) for loss allocation is written as:

$$\begin{aligned} f\{x, y\} &= \sum_{i=1}^{n_L} |loss_{pu}^i - k_{Load_l}^*| \\ &+ \sum_{k=1}^{n_G} |loss_{pu}^k - k_{Gen_l}^*| \end{aligned} \quad (34)$$

where $k_{Load_l}^*$ is the ratio of total system loss to total system load, and $k_{Gen_l}^*$ is the ratio of total system loss to total system generation.

3) *Optimal Tracing Problem Formulation:* Now, OPT problem that was defined in (1) can be explicitly expressed as follows:

$$\min f(x, y) \quad (35)$$

$$\begin{bmatrix} A_d & 0 \\ 0 & A_u \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} b_d \\ b_u \end{bmatrix} \quad (36)$$

$$\begin{aligned} y_k^i P_{L_i} - x_i^k P_{G_k} &\geq 0, \forall k \in \{1, 2, \dots, n_G\} \\ &\forall i \in \{1, 2, \dots, n_L\} \end{aligned} \quad (37)$$

$$[0, 0, \dots, 0]^T \leq [x] \leq [1, 1, \dots, 1]^T \quad [y] \geq [0, 0, \dots, 0]^T \quad (38)$$

where A_d is all the equality constrained matrices corresponding to generation tracing, and A_u is all the equality constrained matrices corresponding to load tracing. By solving above optimization problem we can allocate the generation and load to find the transmission usage cost and loss allocation.

III. SYMBOLIC REDUCTION TECHNIQUE

The tracing problem discussed in Section-II is modeled by a multi-commodity LP problem, where the number of variables are given by following equation:

$$\begin{aligned} No. of Variable &= n_b \times (n_G + n_L) + 2 \times (n_G \\ &\times n_L) + 2(n_G + n_L) \end{aligned} \quad (39)$$

where n_G is number of generators, n_b is number of branches, and n_L is number of loads of the network. The number of variables required for different systems is tabulated in Table I.

Table I No. of variable of different systems

S. No	Name of the system	No. of variables
1	IEEE 6 bus	70
2	IEEE 30 bus	1498
3	IEEE 57 bus	4640
4	IEEE 118 Bus	15252
5	NER Grid (India)	4511420

These numbers are quite large in comparison to the size of the system. In the final LP solution it is seen that many of the variables are set to zero. By identifying this large subset of zeros before the LP formulation we can reduce the size of the problem.

This suggests that additional work is required to improve computational efficiency of the method. This can be achieved by:

1. exploitation of sparsity linear programming and
2. improvement in modeling and algorithm.

The following methodology is proposed for the symbolic reduction. Symbolic reduction approach is a graph theoretic preprocessing on the output of load flow program to identify the reach of generator and load to transmission lines. It also identifies as to which generator can deliver power to loads and vice versa in the tracing frame work. The basic method can be explained as follows:

Let G be the directed graph of the system and arc be represented as $e(o, d)$. If load flow or state estimation output power flow is from node i to node j , then origin o is i and destination d is j . On the other hand if power flow is from j to i then origin o is j and destination d is i . If it is said that the generator at node k can reach the node q then the node q can be reached by traversing a sequence of connected arcs in the direction of flow. In other words, node q is in the reach of node k if following sequence of arcs $e(k, k1), e(k1, k2), \dots, e(kn, q)$ exists. Reach set of generator is the set of all nodes which can be reached by generator. A generator is said to reach an arc $e(l, m)$ if $l \in R_{gk}$ where R_{gk} is the reach set of the generator k . Similarly, the load reach set R_{li} can be found out and load i is said to reach an arc $e(l, m)$ if $m \in R_{li}$.

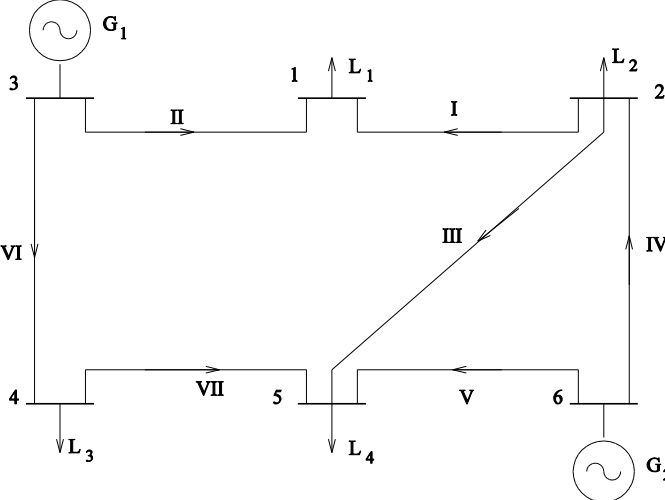


Fig. 2. Graph of a 6 bus system.

This method is illustrated using the Figures 2 and 3. Let us take a directed graph as shown in Figure 2 which is a 6-bus system with 2 generators, 4 loads, and 7 branches. If we start from generator G_1 and traverse in the direction of flow we can reach the nodes 1, 3, 4, and 5 only. So, the power from generator G_1 is flowing through branches II, VI, and VII only and it can supplying to the loads L_1 , L_3 , and L_4 only. This information is tabulated in Table II.

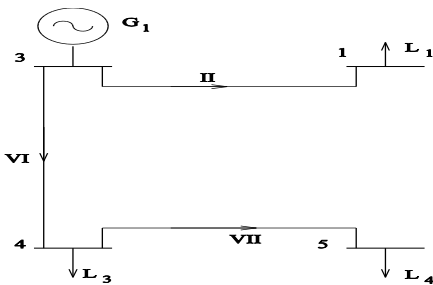


Fig. 3. Reach of Generator G1.
Table II Reach set of generators

Generator	Branches	Loads
G1	II, VI, and VII	L1, L3 and L4
G2	I, III, IV, and V	L1, L2 and L4

Therefore instead of considering all the generators to decompose the flow of a branch, consider only those generators whose power is flowing through that branch. Then the equation (4) is modified as follows:

$$P_{lm} = \sum_{k \in R_{lm}^G} x_{lm}^k \cdot P_{Gk}, \forall \text{ set of lines} \quad (40)$$

where R_{lm}^G generator reach set for line lm i.e., it is the set of all generators which reach line lm .

Similarly, to decompose the load in terms of generator contributions consider the generators from which that load can be supplied. Then the equation (11) is modified as follows:

$$P_{Li} = \sum_{k \in R_{Li}^G} x_i^k \cdot P_{Gk}, \forall i = 1, \dots, n_L \quad (41)$$

where R_{Li}^G generator reach set for load L_i i.e., it is the set of all generators which reach load L_i .

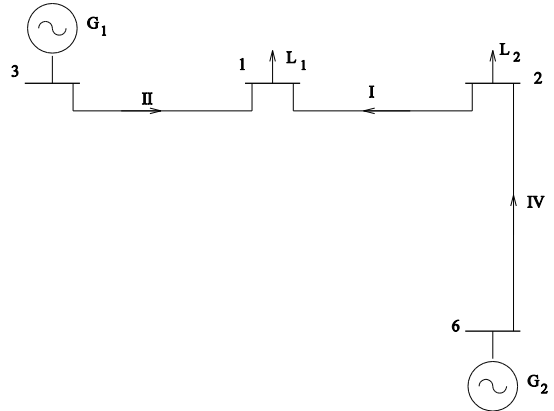


Fig. 4. Reach of load L1.

Table III Reach set of loads

Load	Branches	Generators
L1	I, II, and IV	G1 and G2
L2	IV	G2
L3	VII	G1
L4	III, IV, V, VI, and VII	G1 and G2

From the graph shown in Figure 4, load L_1 is drawing power from both the generators G_1 and G_2 through the branches I, II, and IV (for other loads it is listed in Table III). So, instead of decomposing the branch flow to all load components, we restrict decomposition to only those loads which can reach a line and generator. Thus, equations (6) and (13) can be modified as follows:

$$P_{lm} = \sum_{i \in R_{lm}^L} y_{lm}^i \cdot P_{Li}, \text{ for all set of lines} \quad (42)$$

Where R_{lm}^L load reach set for line lm i.e., it is set of all loads which reach line lm .

$$P_{Gk} = \sum_{i \in R_{Gk}^L} y_k^i \cdot P_{Li}, \text{ for all } k = 1, \dots, n_G \quad (43)$$

where $R_{G_k}^L$ load reach set for generator G_k i.e., it is the set of all loads which reach generator G_k .

In this implementation to traverse the graph and keep the track of visited nodes and edges we used the well-known *depth first search* [6].

For an arc, if a generator or load is not in its reach set, then it implies that the corresponding commodity flow cannot reach the arc. Hence the corresponding generator tracing or load tracing variables can set to zero. This can achieve significant reduction in the number of variables actually modeled by LP. This leads to reduced number of variables in the LP problem. Further, the reduced problem can be solved by sparse LP method.

IV. OBJECT ORIENTED DESIGN [7] FOR EFFICIENT TRACING ENGINE

Implementation of the computationally efficient tracing engine requires sparse linear programs. Any commercial LP package would be sufficed. The tracing engine has mainly two classes one is tracing engine and another one is for core tracing algorithm. The class names are given as follows:

1. TracingEngine
2. TracingAlgo

The class diagram of the above two classes are shown in the Figures 5 and 6.

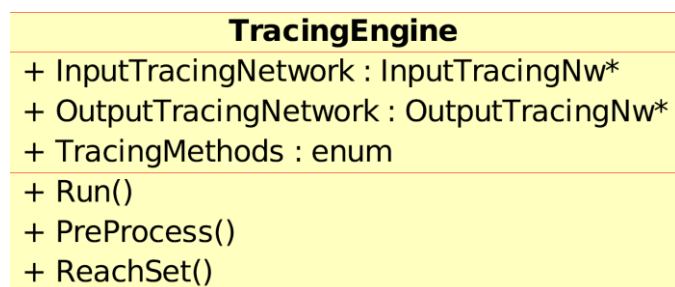


Fig. 5. Class diagrams of main classes of Tracing Engine.

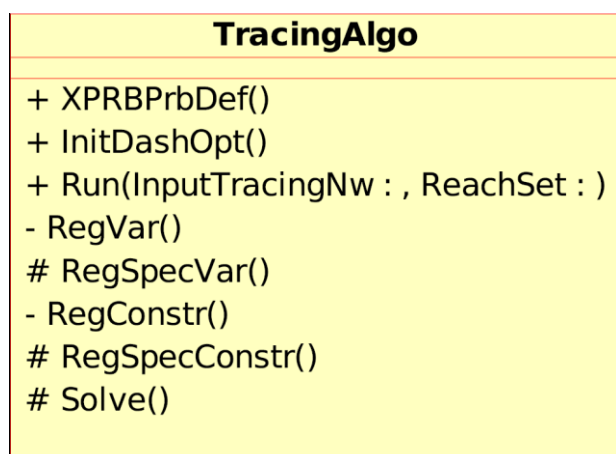


Fig. 6. Class diagram of TracingAlgo.

A. Tracing Engine

The tracing engine is mainly meant for the following tasks.

- build the input
- pre-process the input data
- build input network object
- find the reach set
- compute the output of the tracing engine

B. Tracing Algorithm

1) *Attributes of TracingAlgo class*: The tracing algorithm class (**TracingAlgo**) is a design for core computation of the tracing. This class will call dashoptimization tool box for solving linear programming problem.

2) *Methods in TracingAlgo class*:

- **InitDashOpt()**: Initialize the dash optimization tool.
- **XPRBPrbDef()**: used to define a problem for Dash optimization.
- **Run(InputTracingNw:, Reachset:)**: Run the tracing algo by passing arguments **InputTracingNw** and **ReachSet**.
- **RegVar()**: Register the common variables for all the methods. It is a private method.
- **RegSpecVar()**: Register the specific variables for particular method. It is a protected method.
- **RegConstr()**: Register the common constraints for all the methods. It is a private method.
- **RegSpecConstr()**: Register the specific constraints for particular method. It is a protected method.
- **Solve()**: This method used to call the solve method of dashoptimization

V. RESULTS

The algorithm has been tested on the standard IEEE test systems [8] and a real life Indian Grid system [9].

1. IEEE 6 Bus System
2. IEEE 30 Bus System
3. IEEE 58 Bus System
4. IEEE 118 Bus System
5. NER Grid (India) 488 Bus System [9].

TABLE IV: COMPARISON OF ALGORITHMS WITH AND WITHOUT SYMBOLIC REDUCTION TECHNIQUE

S.No	System	No. of Nodes	No. of Variable		No. of constraints		Time required (Secs)	
			Without reduction	With reduction	Without reduction	With reduction	Without reduction	With reduction
1	IEEE 6 Bus	6	70	41	488	343	0.01	0.007
2	IEEE 30 Bus	30	1498	542	1065	511	0.05	0.03
3	IEEE 57 Bus	58	4640	1131	3293	1190	0.18	0.07
4	IEEE 118 Bus	118	23544	3210	15252	3808	1.4	0.29
5	NER Grid (India)	488	4511420	45703	3995453	72168	#	21.83
# - Could not solve optimal tracing problem without symbolic reduction technique								

The number of variables and the computational time of the algorithm with and without symbolic reduction technique is given in the Table IV. The table clearly shows that symbolic reduction technique significantly reduces number of variables and time required to solve optimal tracing problem. The number of variables for solving optimal tracing problem on NER Grid (India) system with symbolic reduction technique is 45703 and it is solved in 21.83 seconds. However, without application of the symbolic reduction technique, the number of variables is 4511420; a commercial optimization software could not solve the problem. The results of different systems is shown in Table IV. The percentage reduction in variables, constraints, and computational time with symbolic reduction technique is given in Table V and Figure 7.

VI. CONCLUSION

The number of variables and the time required for solving an optimal tracing problem have been reduced significantly with symbolic reduction technique. A commercial optimization tool box which failed to solve the large system like NER grid (India) system without symbolic reduction. However it is solved the same system in 21.83 seconds with symbolic reduction technique.

The symbolic reduction technique is not an addition to the optimal tracing problem, it is a prerequisite for the optimal tracing problem.

Table V Comparison of existing and proposed methods

S.No	System Name	No. of Buses	% Reduction		
			Variables	Constraints	Time
1	6 Bus	6	41.43	29.71	30
2	30 Bus	30	63.82	52.02	40
3	57 Bus	58	75.63	63.86	61.11
4	118 Bus	118	86.37	75.03	79.29
5	NER Grid (India)	488	98.99	98.19	#
#-Could not solve optimal tracing problem without symbolic reduction technique					

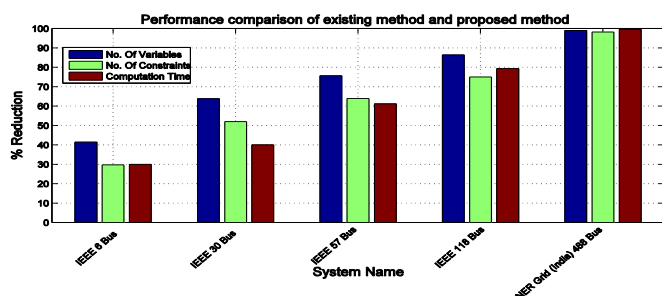


Fig. 7. Percentage reduction in variables, constraints and computational time using proposed method.

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