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# RTOS Partitioning and CWE Compliance: Safeguarding Avionics Software from Critical Failures

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Abstract—In modern avionics systems, Real-Time Operating System (RTOS) partitioning plays a crucial role in ensuring safety, security, and reliability by isolating critical software components. However, software vulnerabilities classified under Common Weakness Enumeration (CWE) can still compromise system integrity if not properly mitigated. This paper explores how CWE rules intersect with RTOS partitioning, emphasizing their impact on flight control, mission-critical data handling, and avionics security. Key vulnerabilities such as buffer overruns (CWE-120), race conditions (CWE-362), deadlocks (CWE-833), and integer overflows (CWE-190) are analyzed in the context of ARINC 653 and other avionics RTOS architectures. We discuss the best practices in inter-partition communication that can prevent these failures, ensuring compliance with DO-178C, ARP4754A, and FACE standards. By integrating CWE-aware development methodologies, avionics engineers can enhance realtime fault tolerance, minimize safety risks, and improve system resilience against cyber threats in partitioned RTOS environments.

## Keywords—Safety, CWE, ARINC 653, DO178B/C

## I. INTRODUCTION (HEADING 1)

In avionics systems, RTOS partitioning is a fundamental technique that ensures isolation between critical software components, preventing faults in one partition from compromising the entire system. However, improper implementation of partitioning can introduce vulnerabilities that lead to catastrophic failures. The Common Weakness Enumeration (CWE) framework provides a structured approach to identifying and mitigating software weaknesses that can arise in partitioned RTOS environments.

Safety-critical avionics software must comply with stringent industry standards such as DO-178C and ARINC 653, where robust partitioning guarantees real-time execution, memory protection, and controlled inter-partition communication. However, buffer overflows (CWE-120), race conditions (CWE-362), and deadlocks (CWE-833) can still emerge within or across partitions if not properly addressed. By integrating CWE-compliant development practices, avionics engineers can ensure fault containment, predictable execution, and high reliability in embedded flight software.

This article explores how RTOS partitioning intersects with CWE rules, highlighting key vulnerabilities and mitigation strategies that enhance avionics software resilience.

#### II. OBJECTIVE

To ensure the safety and reliability of avionics software by implementing RTOS partitioning and CWE compliance, thereby preventing software vulnerabilities such as buffer overruns, race conditions, and deadlocks that could lead to system failures.

Mathematical Expression:

Let:

- S = System reliability
- P = Number of RTOS partitions
- Vi = Probability of a vulnerability in partition iii
- Mi= Mitigation effectiveness for partition iii
- R = Overall risk factor due to software vulnerabilities

The objective is to **minimize risk** and **maximize system reliability**, expressed as:

S=1-R

$$R = \sum_{i=1}^{P} V_i \cdot (1 - M_i)$$

where:

- Vi is reduced by applying CWE rules (e.g., buffer protection, race condition handling).
- Mi is increased by using memory isolation, time partitioning, and deterministic execution in RTOS.

Thus, the goal is to **maximize Mi** for all partitions to achieve  $S \rightarrow 1$  (i.e., a highly reliable system with minimal risk).

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## III. PROBLEM STATEMENT

In modern safety-critical avionics systems, ensuring the reliable operation of software is paramount. RTOS partitioning is a key approach to isolate different software components into distinct, time-controlled partitions to prevent interference and ensure real-time constraints are met. However, the integration of partitioning with CWE (Common Weakness Enumeration) rules is crucial to avoiding critical vulnerabilities that could lead to system failures or safety hazards.

A significant issue arises when incorrect memory management, uncontrolled access to shared resources, or failure to follow time constraints leads to system instability. These weaknesses, such as buffer overflows, deadlocks, and race conditions, can cause catastrophic effects, especially in real-time operating systems (RTOS), where timing and isolation are essential.

## A. Faulty Code in RTOS Partition

Faulty Code in RTOS Partition (Risk of Buffer Overflow) Imagine an avionics system that collects sensor data and stores it in a buffer. The sensor input is then used for control surface adjustments. In this scenario, if there's no check on the data length, a buffer overflow can occur, leading to memory corruption within the RTOS partition.

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
```

## #define MAX\_BUFFER\_SIZE 128

```
void process_sensor_data(const char *sensor_input) {
   char data_buffer[MAX_BUFFER_SIZE];
   strcpy(data_buffer, sensor_input); // CWE-120: Buffer
Overflow, no bounds checking
   printf("Processed data: %s\n", data_buffer);
}
int main() {
   const char *large_input = "This input string is larger than
the buffer and will cause an overflow!";
   process_sensor_data(large_input); // This call causes a
buffer overflow
   return 0;
}
```

#### B. Situational Analysis

Problem with Partitioning and Buffer Overflow:

In a partitioned RTOS like PikeOS, where the Flight Control System (FCS) is isolated from other subsystems (e.g., navigation, communication), the buffer overflow could overwrite adjacent memory belonging to another partition. This can lead to an inter-partition fault or unintended interaction between partitions, violating the isolation principles that RTOS partitioning is designed to enforce.

## C. Methodology and Mitigation

In avionics systems, RTOS partitioning is employed to separate critical and non-critical tasks. The primary goal of partitioning is to ensure that faults or failures in one partition

do not affect the others. The RTOS partitions are designed to have dedicated resources (memory, CPU time, etc.) to prevent unintended interactions between tasks.

- Time partitioning ensures that each task or partition receives deterministic CPU time without interruptions from other partitions.
- Space partitioning isolates memory resources to prevent one task from corrupting or interfering with another.
- Example: In PikeOS or RTEMS, partitions are defined with specific memory addresses and time slots for execution.

The RTOS scheduler ensures that critical tasks, such as Flight Control Systems (FCS) or Navigation functions, receive higher priority over non-critical background tasks, such as log management or diagnostic checks.

- Tasks within each partition are scheduled based on priorities defined according to their safety-critical importance.
- Non-blocking, priority-driven scheduling prevents task starvation, guaranteeing critical functions are always prioritized.
- D. Impact of Partitioning on CWE Mitigation RTOS partitioning isolates tasks, ensuring that:
- Memory Corruption (e.g., buffer overrun) within one partition does not affect other partitions.
- Critical Tasks have exclusive access to time-sensitive resources, ensuring their performance is not interrupted by non-critical tasks.

This isolation reduces the attack surface, making it harder for potential vulnerabilities in non-critical partitions to propagate and impact the core avionics functions.

#### E. Feedback Loop for Continous Improvement

By collecting data on partition failures, CWE rule violations, and runtime anomalies, developers can continuously improve the software architecture and mitigate newly discovered vulnerabilities. This feedback loop leads to:

• Improved Coding Practices: Incorporating lessons learned into future code development.

Enhanced Fault Tolerance: Adapting partition configurations and task schedules to prevent similar failures from occurring.

# IV. MITIGATION STRATEGIES

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- A. Partition Isolation and Boundaries
- RTOS partitioning involves dividing an avionics system into isolated partitions, each running its own application with restricted resources. This prevents the failure of one partition from affecting others, a fundamental principle in ARINC 653 or PikeOS-based systems.
- CWE-119 (Improper Restriction of Operations within the Bounds of a Memory Buffer): Partition isolation ensures that tasks within one partition cannot overrun their allocated memory and corrupt another partition's data.

Formal Verification Approach:

- Model Checking: Formal verification tools like UPPAAL or PRISM can model and verify partition isolation properties.
   By modeling each partition as a finite-state machine (FSM), it can be mathematically verified that a partition's memory bounds will never be exceeded, preventing buffer overruns.
  - Mathematical Algorithm:

```
For each partition i, define:
MemoryRange_i = [start_i,
end i]
```

```
Ensure that for every memory access in partition i:

MemoryAccess_i

MemoryRange i
```

This ensures that memory accesses stay within the bounds allocated to each partition.

- Safety Property:
- Using temporal logic (CTL or LTL), a property can be verified, such as:

```
AG(MemoryAccess_i >= start_i && MemoryAccess i <= end i)
```

This checks that memory access always stays within valid bounds.

#### B. Task Synchronization and Deadlock Prevention

- In avionics systems, multiple tasks may require access to shared resources. Using semaphore-based synchronization or priority-based task scheduling, critical tasks must be executed without interference from non-critical tasks.
- CWE-833 (Deadlock): Partitioning combined with proper synchronization ensures that tasks don't enter deadlock situations.

Formal Verification Approach:

 Algorithmic Approach: Deadlock prevention algorithms like Banker's Algorithm or Resource Allocation Graphs are used to ensure that tasks never enter a deadlock situation by mathematically analyzing resource allocation patterns.

- Mathematical Algorithm:
- Let T be the set of tasks, R the set of resources, and A[i,j] the allocation matrix where i is the task and j is the resource.
- Ensure there is always a safe sequence of task executions:
- SafeSequence(T, R) = {T1, T2, ..., Tn} such that A[i,j] ≤ Available[j] for all i.
- This ensures that tasks are granted resources in a manner that avoids deadlock.
- Formal Verification:

Use formal methods to verify that no circular wait conditions exist:

AG(¬(exists cycle in ResourceAllocationGraph))

This property ensures that the system will never enter a state where tasks are waiting indefinitely for resources.

- C. Preventing Buffer Overrun and Memory Corruption
- CWE-120 (Buffer Overflow): RTOS partitioning restricts each partition's memory access, and buffer overruns are prevented by using bounds-checking techniques like memcpy() with size limits or by using dynamic memory managers that perform overflow checks.

Formal Verification Approach:

- Buffer Bounds Check: Use abstract interpretation or static analysis to ensure that no write operation exceeds the bounds of allocated buffers within a partition.
  - **Mathematical Algorithm:** For a buffer access operation, ensure:

```
For each buffer, Buffer_i
with size N_i, and write
operation w:
if (w >= 0 && w < N_i)
    Access is valid.
else
    Invalid access.</pre>
```

• **Formal Property:** The formal verification model would ensure that for every write to buffer **i**:

```
AG(Write i >= 0 && Write i < Size i)
```

This temporal logic formula ensures that buffer writes are always within the valid range.

- D. Secure Partitioning and Authentication
- CWE-287 (Improper Authentication): In avionics, sensitive data between partitions or systems must be securely authenticated. Partition boundaries prevent unauthorized access, and partition communications are secured via public key infrastructure (PKI) or HMAC.

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TABLE I. IMPACT AND MITIGATION GLIMPSE

## Formal Verification Approach:

- Cryptographic Security Model: Formal methods can be used to verify that all communications across partitions are properly authenticated, ensuring that only authorized partitions can access sensitive data.
  - Mathematical Algorithm: For partition communication, use HMAC for authentication:

HMAC(PartitionKey, Data) =
Expected HMAC

If the computed HMAC matches the expected value, the data is authenticated.

## • Formal Property:

The system model can verify that communication always involves proper authentication:

AG(AuthenticationCheck(Partition\_i) == True)

#### E. Conclusion

In safety-critical avionics systems, RTOS partitioning combined with CWE compliance provides a robust framework for fault isolation, vulnerability detection, and system reliability. By partitioning critical tasks and applying rigorous static analysis during the development process, avionics software developers can create more secure, fault-tolerant, and mission-ready systems.

This methodology emphasizes the importance of structured development practices and rigorous testing in preventing software failures that could jeopardize flight safety.

TABLE I. IMPACT AND MITIGATION GLIMPSE			
CWE ID	Vulnerability	Impact in Avionics	Mitigation Strategies
CIVIE	D 00	RTOS	11 : (6) : 1 (
-120	Buffer Overflow	Can cause	Use snprintf() instead of
-120	Overnow	memory	sprintf(), apply boundary
		corruption, leading to	checks (memcpy_s() or strncpy()), and enable
		crashes or	MPU (Memory Protection
		execution of	Unit) in RTOS to prevent
		malicious	illegal memory access.
		code.	megar memery weeess.
CWE	Out-of-	Leaks	Implement strict boundary
-125	Bounds Read	sensitive data	checks, use safe memory
		or causes	APIs, and enable
		unexpected	MMU/MPU protections
		behavior.	for each RTOS partition.
CWE	Out-of-	May corrupt	Enforce strict RTOS
-787	Bounds Write	memory or	partitioning with memory
		overwrite	boundaries, validate all
		critical	input sizes, and use
		avionics	compiler flags (-fstack- protector or -
		parameters.	D FORTIFY SOURCE=2
			D_TOKTH T_SOURCE 2
CWE	Double Free	Can cause	Implement reference
-415		undefined	counting, avoid manual
		behavior or	memory management in
		system	safety-critical avionics
		instability.	tasks, and use RTOS
			memory pools instead of
CWE	D	C 1 14-	dynamic allocation.
-362	Race Conditions	Can lead to data	Use RTOS mutexes, semaphores, or message
-302	Conditions	corruption in	queues to synchronize
		shared	tasks and prevent
		avionics	concurrent access.
		resources.	
CWE	NULL	Can cause a	Validate all pointers before
-476	Pointer	system crash	dereferencing, use NULL
	Dereference	in avionics	checks, and enable RTOS
		software.	exception handling to
CWE	Hardcoded	Leads to	catch segmentation faults.  Use secure key storage
-798	Credentials	security	mechanisms, avoid
-770	Credentials	breaches and	embedding credentials in
		unauthorized	firmware, and implement
		access to	runtime credential
		avionics	provisioning.
		systems.	
CWE	Server-Side	Can expose	Restrict network access,
-918	Request	avionics	use firewall rules in
	Forgery	communicatio	RTOS, and validate all
	(SSRF)	n channels to	remote requests before processing.
CWE	Improper	cyber threats.  Can lead to	Use stack canaries, control
-119	Improper Restriction of	stack	flow integrity (CFI), and
117	Operations	overflows or	partitioning enforcement
	within	execution of	in RTOS.
	Memory	arbitrary code.	
	Bounds		
CWE	Deserializatio	Can lead to	Use trusted serialization
-502	n of	remote code	libraries, implement
	Untrusted	execution or	whitelisting, and perform
	Data	avionics	signature verification of
		system hijacking.	data.
	<u> </u>	mjacking.	<u> </u>

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