

# Estimation of Content Spread on NDN

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## II. BACKGROUND

This section 2 presents the content distribution systems, Named Data Networking and testbed.

### a) NDN PRIME

**Abstract**—In currently, IP architecture have more important in daily life. But because the IP architecture still has problems. When there are a large number of users Will cause data collisions and Security in data transfer and a lot of costs. Therefore, Named Data Networking or NDN is an interesting new architecture. By using the name as the center, Unlike IP architecture that will use host as the center. NDN has a unique name packet for each package. And forward to the router to find the route to the destination.[1] Naming the data rather than the location enables NDN routers to cache each packet and improve network bandwidth, the better. The spread of content will use a lot of computed bandwidth of NDN to distribute the content with the caching function. this paper will Explain the examination of the spread of content for Name Data Networking by running an implementation of NDN by using CCNx. NDN's content distribution tests have shown that content packet sizes have a greater impact on system workloads and poor CCNx decisions [2] will increase system workloads in the CCNx program. In addition, found That inefficient applications will cause unnecessary calculations to be decoded.

This article, in section1, Explain the experimental content of the initial distribution. section 2, shows the background for content distribution in NDN, for section 3 shows design and experimentation, section 4 describes experimental performance, and sections 5 and 6 show performance improvements and profiling results CCNx.

**Keywords**—Named Data Networking (NDN), Content-Centric Network (CCNx), Content Spread, Performance Examination.

## I. INTRODUCTION

The internet was created with the purpose of providing communication between hosts. the important aspects of content distribution are the host. Lots of requests for the same content must be transferred multiple times over the network from the server, which is clearly not effective.[3] Focusing on the content will have the two main NDN properties, namely 1. Packet with specific name Will be sent according to the search results of the packet name There are two types of packets in the NDN network: interesting packets and data packets. When users need some content will send interesting packets to show requests the data packets are used to respond to interest packets with appropriate content.[4] 2. NDN routers can store a number of forwarded data packet caches. When the interest packet has arrived at the router. The router will have the content requested. Content, which will be sent back to the host immediately instead of forwarding the request to this server. [5] This can reduce a lot of bandwidth in the NDN network and support better content distribution, Therefore, this test will measure the efficiency of the speed of packet content distribution in the active NDN and find a suitable solution to the current bandwidth problem.

From each node of the NDN network will run the daemon program. There is no difference between the client and the router server. But in the network, there are two types of packets which are : (1) Interest Packet and (2) Data Packet in which the packet names are arranged in a hierarchy for flexibility in the operation of the server router. [4]

In NDN's architecture, there are three main components. For forwarding [5] are: (1) Forwarding Information Base (FIB): Responsible for stores forwarding lists that send interest packets to the possible source of matching Data. Unlike IP architecture, that allows more than one list of exported faces (next-hop) for each name prefix. The FIB be populated manually and/or by The NDN routing protocol in the control plane. (2) Pending Interest Table (PIT): Responsible for collecting the unsatisfied interest packets and received faces, so data packets can be returned to nodes interested in data. and (3) Content Store (CS): Responsible used for caching data.

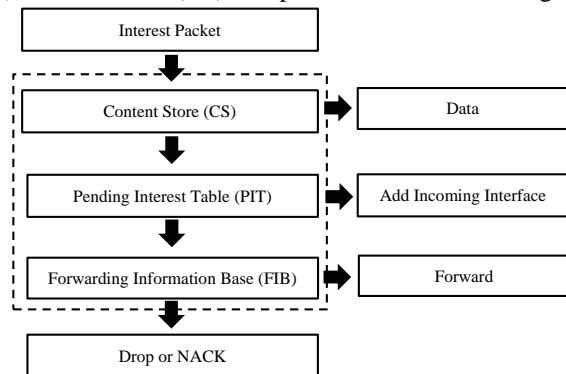


Figure1. forwarding plane of NDN

When attention arrives at the NDN router, the CS will be examined first for the corresponding data. If the CS can respond to the attention, the data packets will be sent back to the receiving face. Otherwise, it will be added to the PIT. If there is an item with the same name, a new face number will be added to the face list to be able to send a copy of the matching data packet on all the faces that the interest packet has arrived. and Eventually, if interest does not have a matching PIT item, it will forward the attention to the next-hop(s) according to the FIB. If there are many more hops in the FIB list, a module called "Forwarding strategies" will determine how to use multiple routes to pass attention on.[6]

b) *Content distribution*

Normally, the distribution of NDN content affects many users. There are many methods, which are Client-Server method the server that stores content and clients can receive content directly from the server to reduce the workload of the server cache. [7] In the second method, a content distribution network (CND) is a network for displaying data and storing in many places. CDN will travel through the network. Gather the closest content with full use of the content delivery area for participants. And the last method is Peer-to-Peer (P2P). In this method it is divided into small sections and the host will download the content according to the group. [8] When the host has finished downloading some content Will continue to act as the content provider for that data the router is close to other preferred hosts, and P2P is still very popular in the network for distributing NDN content and is the easiest way for NDN networks.

c) *Implementation of prototypes*

The forward design of the NDN developed in PARC is a Content-Centric Networking Project (CCNx). The Important content of CCNx is the ccnd daemon.[9] It is a program intended for packets to be forwarding and caching. Generally, NDN does not use IP for real work. If having a real NDN network, IP networks can be used as overlays in the NDN network. In the CCNx project, there are three formats used in the CCNx network are : 1.Data Structures: Logical FIB and PIT use hash tables that provide Name Prefix Hash Table (NPHT), which indexes the Propagating Entries (PEs) and Forwarding Info Entries (FIE) There will be a position to store interest and forwarding information respectively. The Propagating Hash Table (PHT) is sent via the packet space that is shown in PEs and PHT, preventing replay in the network. Each content packet in the CS will receive a unique number. For daemon entry. Content Array (CA) is an array will store cached content packets. Straggler Hash Table (SHT) will collect some of the content packets to increase space efficiency. And there are two types of data structures: Content Hash Table (CHT) and Content Skip List (CSL) for making indicators that store content.

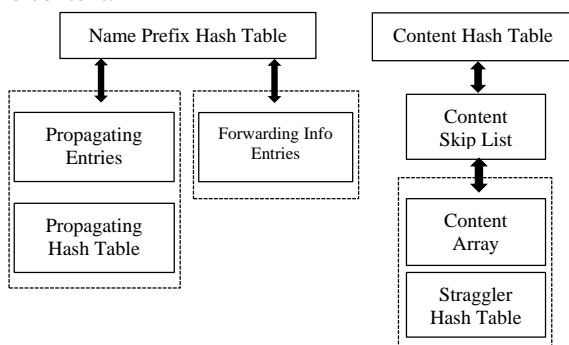


Figure 2. CCNx Forwarding Data Structures

CHT is a hash table with the name of the content packet. and The CSL will arrange a content-order lookup.[6] 2. Interesting packet processing: When the router receives the interest packet. will ask PHT first to make sure that no loops occur. CSL provides with search based on reliable queries and matched content. From CS, content will be pulled out, conducting a comprehensive matching examination. If the

match is satisfactory Content packages will be sent to users who send this attention packet. If the matching is not satisfactory or there is no match, the daemon will find an interesting name on NPHT. The longest prefix match will be performed to find the appropriate FIE.[10] The interest packet will be sent to that host and will be added to PIT and PHT 3. Content packet processing: CHT will ask when the content packet arrives. To ensure that this packet is not stored in CS. If the packet is not repeated, it will be inserted into CSL and CS. After searching on CSL, NPHT will be asked the longest matching prefix to find all interest packets that can be made with this content packet. The content packet will be sent to that host.

III. DESIGN

To assess the content distribution in NDN has to systematically designed our tests and evaluations of the first results obtained. by selecting the metric to use for evaluating the performance of parameters the results obtained from the tests show the effect of the performance in the mentioned systems. and explain the technique for evaluation results and experimental designs 2<sup>k</sup>

a) *System Explanation*

NDN has a distributed system consisting of three components: are NDN Clients, NDN Servers and NDN routers. These three components run the same ccnd daemon program for the corresponding host in the system. To be able to forward and caching packets. The client runs on the client-end software. While the server uses the server-end application running on ccnd.[11] Other applications do not run on the router. But will be forwarded in the content distribution, the content is stored on the server. The client will download the content by sending interest packets. interest packets Respond to application servers by sending content packets. follow the content packet will be cached in the content store of the NDN router.

b) *Indicators*

Examination of the spread of Content for NDN in the context of the NDN. Should choose indicators that show the quality of service the system generates. By using the necessary functions and provided by the content distribution system used on NDN, namely distributing content to end-users.

the ideal indicators are Content Download Time (CDoT) of destination host and Content Distribution Time (CDiT) of the system.[12] CDoT is the time between the client sending the first interest packet and receiving all content packets. CDiT is the time period that sends the first interest packet between clients or all clients that received that content.[13]

For evaluation, NDN router forwarding rate will be used as an evaluation measure. Can measure the NDN forwarding rate using a router based on NP-based. and when the CPU is saturated, the forward rate of the NDN router is set at the maximum throughput that the NDN router can accept for evaluation.

c) *Parameters*

There are many parameters in a computer network that can affect the performance of the NDN-based content distribution system. In this section, all parameters are described before the workload parameters.

Table 1. Pc System Parameters

Parameters	Description
Operating System	Ubuntu16.1
Transmission Protocol	TCP or UDP
Number of Servers	8
Number of Clients	4
Number of Routers	4
Server Application	1
Client Application	1
Network Topology	Tree
Network Reliability	Configurable

System performance studies in Table 1 show software parameters. The CCNx can be configured and modified on each node. While the hardware parameters cannot be modified as easily as the software parameters.

Table 2. Workload Parameters

Parameters	Description
Content Name Length	Configurable
Content Packet Size	Configurable
Content Diversity	Configurable

Since there is no NDN packet tracking, therefore, the design and use of simulated workloads for this evaluation study. Table 2 lists the workload parameters for this performance study.

Table 3. Factors to Study

Factors	Description
Content Store Capacity	1K, 10K, 100K
Content Name Length	2, 4, 8, 16, 32, 64
Content Packet Size	64B, 128B, ..., 1MB
Content Diversity	1, 2, 4, 8, 16

Determinants are selected from the parameters specified in Table1 and Table 2. The main factors of the study listed in Table 3. NDN routers can change the capacity of the configured content store. Usually, CS that are larger can store more content and make it more likely to respond to requests that arrive.[14] The average time of a larger CS increases search and recovery time. Each NDN packet is assigned a unique name. And was assigned the path by name. Many elements of the NDN packet name, each consisting of a slash symbol.

NDN packet name length is determined as a component in the name by Long names increase search time and affect the workload of the system. [15] There is a hypothesis that every host request only one content. The amount of content that is varied is the amount of different content requested by the host.

The parameters are modified using the program. Cncatchunks2 to download the client Which will respond to random data interests Which consists of server number 4 and the client set to 8. This number is set to ensure that the NDN router's CPU can be saturated. In order to be able to measure the maximum workload of the system, choose TCP as the transmission protocol for this assessment. Figure 3 shows the network topology It can be seen that there are 4 NDN servers, namely Server1 to Server4. That are connected to 8 Host NDN is Host1 through Host7 through. A router NDN NP is R1 through R4. Used for network connections and measures the amount of data transmitted by the system.

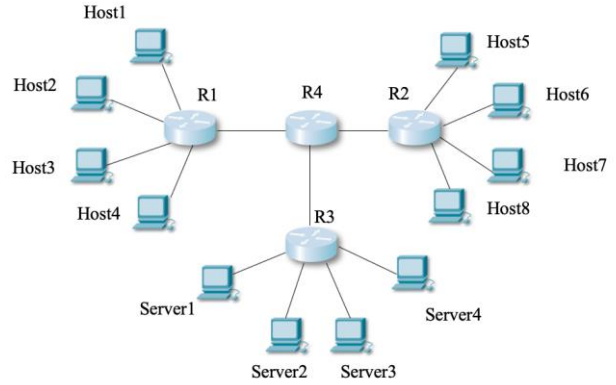


Figure 3. Network Topology

d) Evaluation

The program used in the experiment have been developed and study the measurement of the system using the network topology created using CCNx [16], as have seen in Figure 3, with script development. Some to control the hosts participating in the test the output of the NDN router is measured by a router using NP in Network. The measured data is displayed directly in the Terminal.

e) Experimental Design

Since there are 4 factors in the study Therefore, use the 2<sup>k</sup> factor design for the design of the experiment, where k = 4. This gives the experimental result 2<sup>4</sup> = 16. The list of factors and levels of each factor in Table 4.

Table 4. Factor Symbol and Levels

Symbol	Factor	Level -1	Level +1
A	Content Store Capacity	10K	50K
B	Content Name Length	4	16
C	Content Packet Size	64kB	1MB
D	Content Diversity	1	8

IV. EFFICIENCY

This section provides a description of the performance of the content distribution system on the NDN. Results are presented and explained in this section. By using data analysis techniques to identify key factors affecting distribution efficiency.

Table 5. Measured System Throughput

#	A	B	C	D	Bandwidth Mbps	Package Mbps
1	-1	-1	-1	-1	7.2	41.72
2	1	-1	-1	-1	7.69	45
3	-1	1	-1	-1	8.23	48
4	1	1	-1	-1	8.65	49.19
5	-1	-1	1	-1	9.83	55.78
6	1	-1	1	-1	10.36	58
7	-1	1	1	-1	11	59
8	1	1	1	-1	13.47	64.48
9	-1	-1	-1	1	14.46	67.76
10	1	-1	-1	1	16.62	70
11	-1	1	-1	1	20	75.4
12	1	1	-1	1	22.6	78.7
13	-1	-1	1	1	23.7	89
14	1	-1	1	1	25.27	99.45
15	-1	1	1	1	30	103
16	1	1	1	1	31	121

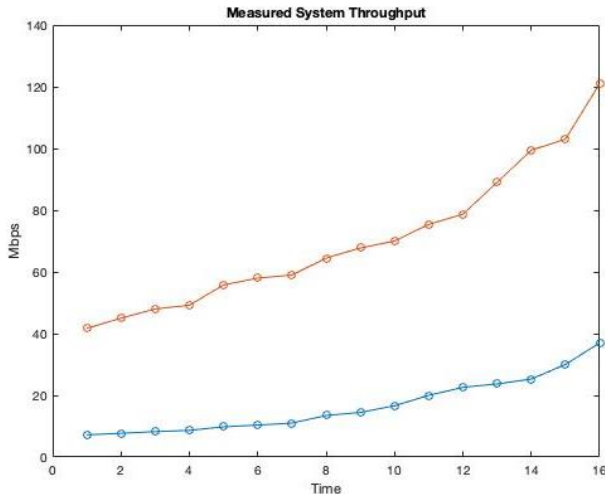


Figure 4 Measured System Throughput

a) *Operating Results*

After the experiment design in Table 5 and Figure 4. There are 16 experiments. The amount of incoming traffic going into the NDN router and the outbound traffic quantity from the router. will be measured using the traffic test tools available on NP-routers. Bandwidth usage is shown in Table. And the value of the workload Can be read directly from SSH. The total workload of the 16 measured systems is in Table 5. With a minimum bandwidth load of 7.2 Mbps and a package starting at 41.72 Mbps in 16 tests, the package volume from the 12th Will increase dramatically, the workload bandwidth will increase step by step, respectively Which the bandwidth and packages are still very far apart.

Table 6. Top 5 Contributors to the Variation

Symbol	Var (%)
C	51
B	22
BC	9.5
D	9
CD	5.6

b) *Data analysis*

Assessing the factors that affect the system and the highest efficiency from the analysis of the data specified in Table 5 or Figure 5, the top 5 factors in terms of participation in the changes specified in Table 6. from the data in Table 6, factors C are Content size and the most important factor affecting the workload of the system. Factor B is the length of the content title, the second priority of the BC factor table is the intersection of two factors. Wash is the third factor. Followed by the D factor is the content diversity, the CD factor is the intersection of the two factors C and D. In addition, Factor A, the storage size of the content is insignificant in terms of impact on system workloads.

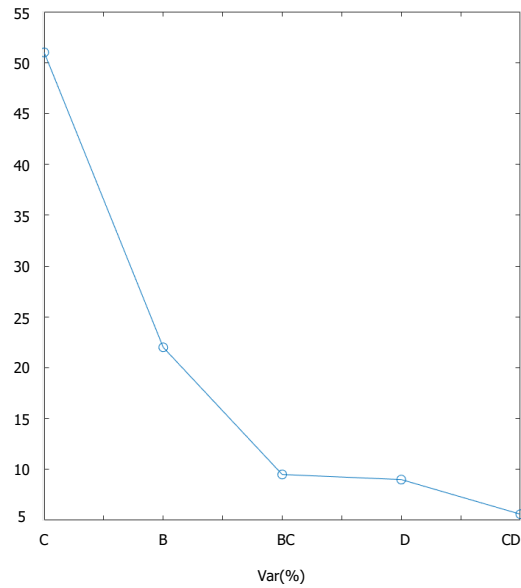


Figure 5. Contributors to the Variation

V. IMPROVING PERFORMANCE

The tests in the previous section had the highest workload used. The speed that can be achieved is 7.2 Mbps and the maximum is 31 Mbps, which results in an increase in volume. But it takes longer for the data to be transmitted more efficiently. Line speed compared to current speed to illustrate that NDN has the ability to provide content distribution functions and other methods of content distribution, which are important in improving and increasing system performance.

This section explains how to increase the forwarding rate. by changing the source code can increase the number of forwarding nodes by up to double to improve forwarding, there are two basic methods: 1) Designing or using data structures for a more effective search for packet names. 2) Rent or buy advanced server hardware to find an efficient route. Today, server hardware has the ability and capability to manage a system on an NDN network. It can find the longest prefix which can increase speed. In addition, the use of nodes for forwarding NDNs on the network processor can improve performance.

Table 7. Top 10 Time-Consuming Functions

Function Name	Percentage (%)
ccn_skeleton_decode	47.34
ccn_buf_advance	7.58
ccn_compare_names	6.1
hashtb_hash	5.87
ccn_buf_match_dtag	3.55
ccn_parse_Signature	3.11
ccn_buf_match_blob	2.01
content_skiplist_findbefore	1.48
hashtb_seek	1.36
content_from_accession	1.21
The Rest	20.39
Sum	100



CCNx is a tool that has the potential to improve the efficiency of the application to suit the current application. the top 10 long-time functions. All functions are from `ccn.h` and `coding.h`, which show the CPU processing time of decoding in Table 7 and Figure 6-4 shows the italic function as a function. Related to decoding packets, as in Table 7, more than half of the time it takes to decipher a particular packet name. `ccn_skeleton_decode` That takes as much as 47.34 percent.

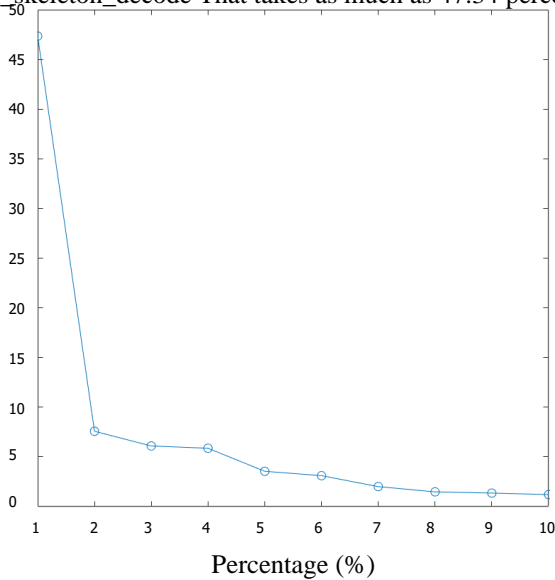


Figure 6. 10 Time-Consuming Functions

Which takes an unusually long time in general network programs. What causes the problem is the content packets to be stored in the Content Store (CS) are encoded. When inquiring about skipping content in the log (n) packet, it must be decoded immediately. Where n is defined as the number of items stored in the Content Skip List (CSL). Which will take a very long time When making program changes. In which the name of the content packet that is decoded Will be stored in the daemon. In this way, there will be no special packet decoding when CSL is asked each time.

## VI. CONCLUSION

In this article, assess the effectiveness of the content distribution system for the named data network. Which will check the workload of NDN. The experiment will show that the workload of CCNx in the operating system Still far from current demand and experiments show that the key factor is the size of the content packet may have the greatest impact on performance and the number of storage name prefixes that will affect. Forwarding if it is too small in comparison to the current application.

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