Duality Gap in Nonlinear Programming

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Abstract:
Duality has an important role in nonlinear programming. It gives a numerical and theoretical foundation for many optimization algorithms. Duality method can be used to solve NLP directly or indirectly as well as it is also useful for finding the upper or lower bound of objective function within its constraints. But duality theory has some limitations. When we use duality for convex problem, then it is best suited. But when we apply duality to non convex problem including discrete and mixed inter problems. It is not always easy to prove weak duality, strong duality, strict converse duality and converse duality theorems and zero duality gaps. In this paper, I have proposed an enhanced duality theory for nonlinear optimization in order to overcome from some limitations of previous dual methods. This enhanced duality theory leads to zero duality gap for nonlinear programming problem which was not possible through previous available methods. and in the last, I have discussed an example which proves that proposed enhanced duality theory is more efficient and effective than others available methods.

Keywords: Nonlinear Programming, Non-convex optimization, duality gap
1. Introduction:

The nonlinear programming problem (NLP) of the following form:

\[ \text{NLP}: \quad \text{Min } f(u) \]
\[ \text{Subject to} \]
\[ l(u) = 0, \]
\[ g(u) \leq 0 \]  \hspace{1cm} (1)

Where variable \( u = (x, y), x \in X \) is the continuous part, and \( y \in Y \) is discrete part, when \( X \subset \mathbb{R}^n \), \( Y \) is finite discrete set of \( k \)-element. We assume that the objective function \( f \) is lower bounded and is continuous and differentiable with respect to \( x \), where \( g = (g_1, g_2, g_3, \ldots, g_l)^T \) and \( h = (h_1, h_2, h_3, \ldots, h_m)^T \) are continuous in the continuous subspace \( X \) for any \( y \in Y \). The NLP covers all type of nonlinear optimization problems.

Duality is an important development for mathematical programming. An important issue of duality theory is the existence of duality gap, which is the difference between the optimal solution of the original problem and the lower bound obtained by solving the dual problem. The duality gap is zero for convex optimization problem generally, but in case of non-convex problem, it is not always true. The duality gap is generally nonzero for non-convex problems and may be large for some problems, in which case the dual approach is not useful. More-over, for discrete and mixed inter problems, duality gap may be nonzero even If the functions are convex. None zero gaps make the direct dual methods fail and the global optimization algorithms such as branch and bound less effective. There is number of research paper [3, 17] are available that has given a efficient method for
removing duality gap in case non-convex problem and due to some limitations in previous works motivated much more for extensive study on this topic.

I have found that some previous dual functions require large penalty values to ensure zero duality gaps for non convex problems. The large penalty values gives ill conditioned optimization problems and increases the difficulty for obtaining optimal solution. Since some previous zero gap results are developed for continuous and differentiable problems. Often the parameters for achieving zero-duality are related the Lagrange’s multipliers [19, 23] do not exist for discrete problems. Such result can be applied to discrete or mixed integer’s problems or problem with nondifferentiable constraints. However I have developed a duality theory in this research for the efficiency and effectiveness on duality gaps.

2. Pre-requisites and related pervious work:
Duality in nonlinear programming or for any mathematical programming is, generally, speaking, the statement of a relationship of a certain kind between two mathematical programming problems. The relationship commonly has three aspects.

One problem the “Primal” is a constraint optimization problem.

The existence of a solution to one of these problems ensures the existence of a solution of the other, in which case their respective extreme values are equal and

If the constraints of the one problem are consistent while those of the other are not, there is a sequence of points satisfying the constraint of the first on which its objective function tends to infinity.

Consider the NLP as:

\[(NLP): \quad \text{Maximize } f(u)\]
Subject to $U = \{ u : u \in U_0, g(u) \leq 0. \}$

Where $U_0$ is an open subset to $\mathbb{R}^n$; $f : U_0 \rightarrow \mathbb{R}$ and $g : U_0 \rightarrow \mathbb{R}$ are differentiable.

Wolfe type dual [ ]:

(WD):

\[ \text{Minimize} \quad f(z) + r^T g(z), \]
\[ \text{Subject to} \quad \nabla f(z) + \nabla r^T g(z) = 0, \]
\[ r \geq 0. \]

Wolfe established weak duality under the assumption that $f$ and $g$ are convex at $U_0$.

Mond and Weir type dual [18].

In relation to (NLP) Mond and Weir [ ] established its corresponding dual as:

(MWD):

\[ \text{Minimize} \quad f(z) \]
\[ \text{Subject to} \quad \nabla f(z) + \nabla r^T g(z) = 0, \]
\[ r^T g(z) \geq 0. \]
\[ r \geq 0. \]

Mond and weir established the weak duality theorem under assumption that $f$ is pseudo convex and $r^T g(z)$ is quasi-convex on $U_0$.

Since duality is an important notion for mathematical programming. Let consider the following continuous nonlinear programming (CNLP) problem.

(CP):

\[ \min f(x) \quad \text{where} \quad x = (x_1, x_2, \ldots, x_n)^T \in X \]

\[ \text{Subject to} \]
\[ h = (h_1, h_2, h_3, \ldots, h_m)^T = 0 \quad \text{and} \quad (2) \]
\[ g = (g_1, g_2, g_3, \ldots, g_t)^T \leq 0. \]

X is a compact subset of \( \mathbb{R}^n \), f, g, and h are lower bounded and continuous, but not necessarily differentiable.

Then the Lagrangian function of the form:

\[ L(x, \gamma, \delta) = f(x) + \sum_{i=1}^{m} \gamma_i h_i(x) + \sum_{j=1}^{t} \delta_i g_j(x) \quad (3) \]

Dual methods transform the original problem into a dual problem defined as follows.

\[ \text{(CP dual)} \quad \text{maximize} \quad K(\gamma, \delta) \]

\[ \text{Subject to} \quad \gamma \in \mathbb{R}^m \text{ and } \delta \geq 0. \quad (4) \]

Where the dual function \( K(\gamma, \delta) \) is defined as:

\[ K(\gamma, \delta) = \inf L(x, \gamma, \delta) = \inf \left[ f(x) + \sum_{i=1}^{m} \gamma_i h_i(x) + \sum_{j=1}^{t} \delta_i g_j(x) \right] \quad (5) \]

The main results of the duality theory are the following. First, the objective value \( L \) obtained from solving the dual problem (CP dual) is a lower bound to the optimal objective value f of the original problem. There have been a number of studies for elimination the duality gaps. A number of previous works have indicated that duality gap can be reduced when a problem is decomposed or has certain special structures [1.2.21].

To remove the duality gaps for nonconvex problems, augmented Lagrangian functions [3, 17] were introduced for continuous NLP. Rubinov et al. [19, 23] have extended the \( l_1 \) penalty function to a class of nonlinear penalty function with zero duality gaps, where the function takes the following form.

\[ l_\gamma(x, c) = \left[ f^*(x) + c \left( \sum_{i=1}^{m} I_{i} h_i(x) I + \sum_{j=1}^{t} g_j^+(x) \right) \right]^{1/\gamma} \quad (6) \]
Where $\gamma > 0$ is a parameter.

Luo et al. [14] have proposed a nonconvex and non-smooth penalty function with zero duality gap based on the following formulation, where $\gamma > 0$ is a parameter.

$$l_\gamma(x, c) = f(x) + \left[c \left(\sum_{i=1}^{m} I \ h_i(x)I + \sum_{j=1}^{r} g_j^+(x)\right)\right]^{\gamma} \quad (7)$$

An exact penalty function with zero duality gaps under certain assumptions is proposed by Pang [16] as follows.

$$l_\gamma(x, c) = f(x) + c\left[\max\{ I \ h_1(x) I, I \ h_2(x) I \ldots I \ h_m(x) I, I g_1^+(x) I, \ldots I g_t^+(x) I \} \right]^{\gamma} \quad (8)$$

In case of continuous problem as given in (2), most of the existing augmented Lagrangian functions and exact penalty functions with zero duality gap for nonconvex continuous optimization [5-7, 14-15, 18-20, 23].

$$l(x, \gamma, \delta, c) = f(x) + \tau(\delta, \gamma, h, g) + c\sigma(h, g) \quad (9)$$

Where $\delta, \gamma$ are Lagrange multiplier vector, $\tau(\delta, \gamma, h, g)$ is Nonlinear Lagrangian terms, $c \geq 0$ is a penalty parameter and $c\sigma(h, g)$ is an augmenting function. A general framework that provides a unified treatment for a family of Lagrange type functions and conditions for achieving zero duality gap for constrained optimization problems under some convexity assumptions is given by Burachik and Robonov [1, 5]. A recent work by Nedic and Ozdaglar [15] develops necessary and sufficient conditions for $l(x, \gamma, \delta, c)$ to have zero duality gap based on a geometric analysis, which consider the geometric primal problem of finding the minimum intercept. The most previous methods for removing the duality gaps use a single penalty multiplier $c$ before augmenting term. However a best $c$ is hard to
locate and control. In practice a common problem is that the single c is often too large, this makes the unconstrained optimization difficult. In this paper I have proposed multiple penalty multipliers which effectively reduce penalty values for ensuring zero duality gaps.

3. Proposed work for continuous nonlinear programming problem:
I have develop our results for continuous nonlinear programming problems (CPs) (2) given as below.

Definition: 3.1 (Constrained global minimum):
A point $x^*$ in X is a constrained global minimum, if $x^*$ is feasible and $f(x^*) \leq f(x)$ for all feasible x in X.

Definition: 3.2 The $l_1^m$ penalty function for (CP) in (2) is defined as:
$L_m(x, \alpha, \beta) = f(x) + \alpha^T I h(x) I + \beta^T g^+(x)$  \hspace{1cm} (10)
Where
$|h(x)| = (|h_1(x)|, \ldots, |h_m(x)|)^T$ and $|g^+(x)| = (|g_1^+(x)|, \ldots, |g_i^+(x)|)^T$
Where we defined $\phi^+(x) = \max(0, \phi(x))$ for a function $\phi$, and $\alpha \in R^m$ and $\beta \in R^l$ are penalty multipliers.

Definition: 3.3
The directional derivative of a function $f : R^n \rightarrow R$ at a point $x \in R^n$ along a direction $p \in R^n$ is
$f'(x; p) = \lim_{\epsilon \rightarrow 0} \frac{f(x+\epsilon p)-f(x)}{\epsilon}$  \hspace{1cm} (11)

Definition: 3.4 (Constraint Qualification condition)
A point $x$ in $X$ of (CP) satisfies the constraint qualification if there exist no direction $p \in \mathbb{R}^n$ along which the directional derivatives of the objective is non-zero but the directional derivatives of continuous equality and continuous active inequality constraints are all zero. That is, there does not exist $p \in \mathbb{R}^n$ such that

$$f'(x; p) \neq 0, h'(x; p) = 0 \text { and } g_j'(x; p) = 0, \forall \in C_h \text { and } C_g$$

Respectively, the sets of indices of continuous equality and continuous active inequality constraints. The constraint qualification is satisfied if both $C_h$ and $C_g$ are empty. Intuitively, constraint qualification at $x$ ensures the existence of finite $\alpha$ and $\beta$ that lead to a local minimum of (10) at $x$. Consider a neighboring point $x + p$ infinitely close to $x$, where the objective function $f$ at $x$ decreases along $p$ and all active constraints at $x$ have zero directional derivative along $p$. In this case, all the active constraints at $x + p$ are close to zero, and it will be impossible to find finite $\alpha$ and $\beta$ in order to establish a local minimum of (10) at $x$ with respect to $x + p$. To ensure a local minimum of (10) at $x$, the above scenario must not be true for any $p$ at $x$.

We compare our constraint-qualification condition to the well-known regularity condition in KKT condition that requires the linear independence of gradients of active constraint functions. Of course, the regularity condition only applies to differentiable problems, while our constraint-qualification condition does not require.

**Definition 3.5 (Feasible set and $\epsilon$-extension):**

Let the set of all feasible points of (CP) be:

$$F = \{ x \mid x \in X, h(x) = 0, g(x) \leq 0 \}, \ (12)$$

the $\epsilon$-extension of $F$, where $\epsilon > 0$ is a scalar value, is:

$$F^+ = \{ x \mid x \in X, \min ( \| y - x \|) \leq \epsilon \} \ (12).$$

Namely, $F^{\epsilon +}$ includes the points in $F$ and all those points whose projection distance
to F is within $\epsilon$. Here, II, II denotes the Euclidean norm.

**Preposition 3.1:** If $f(x)$, $h(x)$ and $g(x)$ in (CP) are differentiable, if a constraint global minimum $x$ in $X$ of (CP) is regular, then it satisfies the constraint qualification.

**Theorem 3.1 (Enhanced duality theorem for continuous programming):**

Suppose $x^* \in X$ is a Constrained global minimum to (CP) and $x^*$ satisfies the constraint qualification, then there is no duality gap for the enhanced dual problem defined in (10) and (11), i.e. $q^* = f(x^*)$.

**Proof:** First, we have $q^* \leq f(x^*)$. Since

$$q^* = \max_{\alpha \geq 0, \beta \geq 0} q(\alpha, \beta) = \max_{\alpha \geq 0, \beta \geq 0} (\min_{x \in X} L_m(x, \alpha, \beta))$$

$$\leq \max_{\alpha \geq 0, \beta \geq 0} (\min_{x \in X} L_m(x^*, \alpha, \beta)) = \max_{\alpha \geq 0, \beta \geq 0} f(x^*) = f(x^*) \quad (13)$$

Also according to theorem 3.1, there are $\alpha^{**} \geq 0$ and $\beta^{**} \geq 0$ such that

$q(\alpha^{**}, \beta^{**}) = f(x^*)$, we have

$$q^* = \max_{\alpha \geq 0, \beta \geq 0} q(\alpha, \beta) \geq q(\alpha^{**}, \beta^{**}) = f(x^*) \quad (14)$$

Since $q^* \leq f(x^*)$ and $q^* \geq f(x^*)$, we have $q^* = f(x^*)$.

Taking the following DNLP

(Pd): $\min f(y)$, Where $y = (y_1, \ldots, y_w)^T \in Y$

Subject to $h(y) = 0$ and $g(y) \leq 0$ \quad (15)

Whose $f$ is lower bounded, $Y$ is a finite discrete set, and $f$, $g$ and $h$ are not necessarily continuous and differentiable with respect to $y$.

**Definition 3.6 (Constrained global minimum of Pd):**
A point \( y^* \in Y \) is a CGM, a constrained global minimum of \( P_d \), if \( y \) is feasible and \( f(y^*) \leq f(y) \) for all feasible \( y \in Y \).

**Theorem 3.2:** Let \( y^* \in Y \) be a constraint global minimum to (DP), there exist finite \( \alpha^* \geq 0 \) and \( \beta^* \geq 0 \) such that

\[
f(y^*) = \min_{y \in Y} L_m(y, \alpha**, \beta**) \text{, for any } \alpha** \geq \alpha^*, \beta** \geq \beta^*. \tag{16}
\]

**Proof:**

Given \( y^* \), since \( L_m(y, \alpha**, \beta**) = f(y^*) \) for any \( \alpha^* \geq 0 \) and \( \beta^* \geq 0 \), we need to prove that there exist finite \( \alpha^* \geq 0 \) and \( \beta^* \geq 0 \) such that

\[
f(y^*) \leq L_m(y, \alpha**, \beta**) \text{, for any } \alpha** \geq \alpha^*, \beta** \geq \beta^* \text{ for any } y \in Y. \tag{17}
\]

We take the \( \alpha^* \) and \( \beta^* \) such that:

\[
\alpha^*_i = \max_{y \in Y, |h_i(y)|} \left\{ \frac{f(y^*) - f(y)}{|h_i(y)|} \right\}, i = 1, 2...m \tag{18}
\]

\[
\beta^*_j = \max_{y \in Y, g_j(y) > 0} \left\{ \frac{f(y^*) - f(y)}{g_j(y)} \right\}, j = 1, 2...t \tag{19}
\]

Next, we show that \( f(y^*) \leq L_m(y, \alpha**, \beta**) \text{ for any } y \in Y, \alpha** \geq \alpha^*, \beta** \geq \beta^* \)

For a feasible point \( y \in Y \), since \( h(y) = 0 \) and \( g(y) \leq 0 \),

We have \( L_m(y, \alpha**, \beta**) = f(y) \geq f(y^*) \). \tag{20}

For an infeasible point \( y \in Y \), if there is at least one equality constraint \( h_i(y) \) that is not satisfied (\( |h_i(y)| > 0 \)), we have:

\[
L_m(y, \alpha**, \beta**) = f(y) + \sum_{i=1}^{m} \alpha^*_i |h_i(y)| + \sum_{j=1}^{t} \beta^*_j g_j^+(y) \geq f(y) + \alpha^*_i |h_i(y)| \geq f(y) + \frac{f(y^*) - f(y)}{|h_i(y)|} |h_i(y)| = f(y^*) \tag{21}
\]

If there is at least one inequality constraint \( g_j(y) \) that is not satisfied (\( g_j(y) > 0 \)),

We have:
Equation (16) is proved after combining (20), (21) and (22).

The extended dual problem for Pd is the same in definition 3.2 defined for CP, except that the variable space is Y instead of X. Based on Theorem 3.3, we have the following result for discrete-space extended duality, which can be proved in the same way as the proof to Theorem 3.2.

**Example:**
I have illustrated a continuous problem where there is a duality gap for the original duality theory but not for the proposed enhanced extended duality theory. Consider the following CNLP:

\[
\begin{align*}
\min_{x \in \mathbb{R}^2} & \quad f(x) = x_1 + x_2 \\
\text{Subject to} & \quad h(x) = x_1 x_2 - 1 
\end{align*}
\]

It is obvious that \( f^* = 2 \) at \( (x_1^*, x_2^*) = (1, 1) \), for the original duality, the dual function is:

\[
q(\mu) = \min_{x \geq 0} L(x_1, x_2, \mu) = \min_{x \geq 0} L(x_1 + x_2 + \mu(x_1 x_2 - 1))
\]

We have \( q^* = \max_{\mu \in \mathbb{R}} q(\mu) \leq 0 \). As result, there is a nonzero gap since \( f^* - q^* \geq 2 - 0 = 2 \).

In contrast using proposed duality theory, the extended dual function is:

\[
q(\alpha) = \min_{x \geq 0} L(x_1, x_2, \alpha) = \min_{x \geq 0} L(x_1 + x_2 + \alpha(x_1 x_2 - 1))
\]

It is easy to validate that, for \( \alpha ** \geq \alpha ** = 2, q(\alpha) = \min_{x \geq 0} L(x_1, x_2, \alpha) = 2 \).

Therefore, we have \( q^* = f^* = 2 \) and there is no duality gap for the proposed duality approach.

**4. Conclusions**
In this paper, we have proposed the theory of duality for nonlinear optimization. The theory overcomes the limitations of conventional duality theory by providing a duality condition that leads to zero duality gap for general nonconvex optimization problems in discrete, continuous and mixed spaces. Based on proposed penalty function, the proposed theory requires less penalty values to achieve zero duality gap comparing to previous efforts for removing the duality gap, thus alleviating the ill conditioning of dual functions.

**References:**


