

Dependable and Secure Distributed Data Storage in Cloud Computing

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Abstract

Cloud computing is a type of computing that provides simple, on-demand access to pools of highly elastic computing resources. These resources are provided as a service over a network (often the Internet), and are now possible due to a series of innovations across computing technologies, operations, and business models. Cloud enables the consumers of the technology to think of computing as effectively limitless, of minimal cost, and reliable, as well as not be concerned about how it is constructed, how it works, who operates it, or where it is located. Cloud computing is a style of computing where computing resources are easy to obtain and access, simple to use, cheap, and just work. Cloud data user does not possess direct control of his data, security is one of the few challenging issues which need to be addressed. Security in Cloud Computing can be addressed in many ways viz. authentication, integrity or correctness, confidentiality and data error localization. Data integrity or correctness is an issue where there may be some unauthorized alteration in the data without consent of the data owner. To prevent data access from unauthorized access, it proposes a distributed scheme to provide security of the data in cloud. This could be achieved by using homomorphism token with distributed verification of erasure-coded data. Proposed scheme perfectly stores the data and identifies the any tamper at the cloud server and also performs some operations like data updating, deleting and appending. The proposed design allows users to audit the cloud storage with very lightweight communication and computation cost.

Keywords- Cloud computing, Data integrity, Cloud service provider (CSP), Homomorphic encryption and Dynamic operations.

1. Introduction

Many companies today are expanding into cloud computing as a way to reduce the cost and complexity of delivering traditional IT services. Cloud is not a particular product, but a way of delivering IT services that are consumable on demand, elastic to scale up and down as needed, and follow a pay-for-usage model.

Cloud is an elastic delivery model that enables businesses to become more adaptable and interconnected. Monolithic and ageing infrastructures give way or progress toward a 'rent versus buy' state of agility, where non-core competencies are shed for not just on-demand technology, but also for on-demand business innovation and savings. Cloud is not a point product or a singular technology, but a way to deliver IT resources in a manner that provides self-service, on-demand and pay-per-use consumption.

Cloud is designed to distribute IT resources in a cost-effective and nimble way. Consumption-driven cloud commerce moves an organization's focus from CAPEX (capital expenditure), which typically isn't fully utilized, to smaller, incremental and variable OPEX (operating expenditure).

The US Department of Commerce's National Institute of Standards and Technology (NIST) defines Cloud Computing as: "a model for enabling ubiquitous, convenient, on-demand network access to a shared pool of configurable computing resources (e.g., networks, servers, storage, applications, and services) that can be rapidly provisioned and released with minimal management effort or service provider interaction." [1] CLOUD (Common Location independent Online Utility on Demand) is a broad solution that delivers IT as a service. Cloud computing is an umbrella term used to refer to Internet based development and services.

Cloud Computing [2] is a general term used to describe a new class of network based computing that takes place over the Internet. Cloud computing shared resources are provided like electricity distributed on the electricity grid.

A cloud platform service provider (CPSP, e.g., Amazon.com, Google.com, Salesforce.com, and others) provide cloud-based platforms, hosted in a cloud-enabled infrastructure and cloud operating system environment, such that developers can access the platform, develop a new business application, and then host that application on the cloud-based platform. Cloud platform service providers are unique in that they have developed a complete application platform, hosted in a cloud, which enables rapid application development

on that platform, while providing an “as a Service” deployment and hosting framework for the applications to be provided “as a Service” through that platform, which is in turn hosted on a cloud.

2. Proposed Scheme

2.1. Proposed System Architecture

Fig. [3] is architecture representing proposed system which provides dependable and secure distributed data storage service in cloud computing. Admin module proposed in this architecture reduces overhead of users of generating keys.

Three different main components can be identified as follows:

2.1.1. User: an entity, who has data to be stored in the cloud and relies on the cloud for data storage and computation, can be either enterprise or individual customers.

2.1.2. Cloud Server (CS): an entity, which is managed by cloud service provider (CSP) to provide data storage service and has significant storage space and computation resources (we will not differentiate CS and CSP hereafter).

2.1.3. Third-Party Auditor: an optional TPA, who has expertise and capabilities that users may not have, is trusted to assess and expose risk of cloud storage services on behalf of the users upon request.

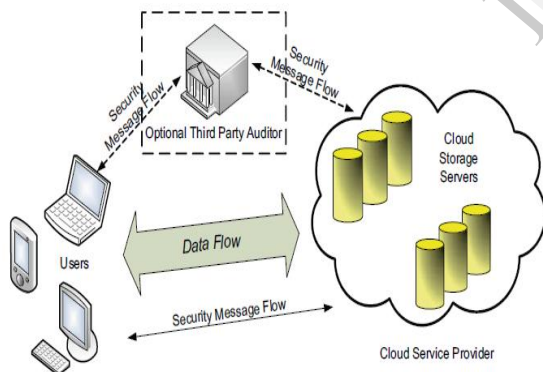


Fig.1. Cloud storage service architecture[3]

In the proposed system architecture, data which is uploaded by users is divided into multiple blocks and stored across selected cloud servers but not randomly across servers as specified in other related schemes. In this paper, as per the proposed scheme, cloud servers are selected based on constraints like cost and quality. Thereby ensuring efficient cloud data storage.

The Admin module is responsible for ensuring authorized access of data stored across cloud servers by restricting unauthorized access of cloud data,

also responsible for generation of master key used by users for creation of digital signature and also generates public key used by auditors during their auditing scheme and thereby reduces overhead of user module of generating keys.

In cloud data storage, a user stores his data through a CSP into a set of cloud servers, which are running in a simultaneous, cooperated, and distributed manner.

- Data redundancy can be employed with a technique of erasure-correcting code to further tolerate faults or server crash as user's data grow in size and importance.
- Thereafter, for application purposes, the user interacts with the cloud servers via CSP to access or retrieve his data.

2.2. Adversary Model

From user's perspective, the adversary model has to capture all kinds of threats toward his cloud data integrity. Because cloud data do not reside at user's local site but at CSP's address domain, these threats can come from two different sources: internal and external attacks.

For internal attacks, a CSP can be self-interested, untrusted, and possibly malicious. Not only does it desire to move data that has not been or is rarely accessed to a lower tier of storage than agreed for monetary reasons, but it may also attempt to hide a data loss incident due to management errors.

For external attacks, data integrity threats may come from outsiders who are beyond the control domain of CSP, for example, the economically motivated attackers. They may compromise a number of cloud data storage servers in different time intervals and subsequently be able to modify or delete users' data while remaining undetected by CSP.

Therefore, we consider the adversary in our model has the following capabilities, which captures both external and internal threats toward the cloud data integrity.

2.3. Design Goals

Our main goal is to ensure the security and dependability for cloud data storage and to design efficient mechanisms for dynamic data verification and operation as follows:

- Pre-computation token key generation algorithm which is simple, elegant and secure method and less overhead due to few parameters that has to be chosen.

- Challenge verification scheme was designed in easy and efficient way to prevent data from Byzantine server failures and data dependability detection or detect data errors on blocks.
- Cloud servers ensure that the file was saved successfully without block modifications. This can be achieved by two way token checking.

3. Ensuring Distributed Data Storage over Cloud

In cloud distributed data storage system, users' store their data remotely i.e., on clouds, so that the correctness and availability of data files being stored on the distributed cloud servers must be guaranteed. Our main aim is to detect the servers which behaves differently and may leads to internal and external threats.

In this paper, we explore the techniques used to detect the modified blocks easily with very less overhead using homomorphic token pre-computation, correctness verification and error localization and file retrieval and error recovery techniques to acquire the desired blocks from different servers.

3.1. Challenge Token Pre-computation

To achieve distributed data storage correctness and data integrity, we use an algorithm which takes a few parameters and compute the token. Here we assume Third Party Auditor (TPA) will participate in key generation [4].

Token generation algorithm works as follows: Let 'f' be the filename and 'fl' be the length of the file and v be the secret matrix which contains special characters in randomized order.

Compute the key with the following parameters:

Algorithm-1. Token Precomputation

procedure

Choose parameters f, fl and secret vector v

Choose number of blocks to be taken (normally fixed block size)

$X = f + fl + v$

Compute key

for $i = 1$ to n

$filetoken = filetoken + (\sum_{i=1}^n split(X_i))$

Compute short signatures for each block of the file by considering token and file block data using bit permutations (token+block data) and store these values in client for dynamic checking

end procedure

Before file is distributed to the cloud, TPA will generate token key with required parameters passed by user, once the token key has been

generated, TPA will send the file by dividing the file into equal sized blocks and generate a small token signature for each block along with initial key filetoken. This filetoken was generated based on mathematical calculations with hash based technique, it is fully randomized we are not explore the operations present in it and just given the function $split(X)$.

Before sending the block it stores the computed signatures obtained from bit permutations on both filetoken and block data. The resultant token was stored in its database or at clients place. Each block is send along with short signature and each block is treated as encrypted block.

3.2. Correctness Verification and Error Localization

Error localization is a key prerequisite for eliminating errors in storage systems. It is also of critical importance to identify potential threats from external attacks. Our proposed scheme outperforms those by integrating the correctness verification and error localization (misbehaving server identification) in our challenge-response protocol: the response values from servers for each challenge not only determine the correctness of the distributed storage, but also contain information to locate potential data error(s).

Once the inconsistency among the storage has been successfully detected, we can rely on the precomputed verification tokens to further determine where the potential data error(s) lies in. Note that each response $R_i^{(j)}$ is computed exactly in the same way as token $v_i^{(j)}$, thus the user can simply find which server is misbehaving by verifying the following n

$$R_i^{(j)} \stackrel{?}{=} v_i^{(j)}, j \in \{1, \dots, n\}.$$

equations:

Algorithm 2. Correctness Verification and Error Localization [3].

procedure CHALLENGE(i)

Recompute $\alpha_i = f_{k_{chal}}(i)$ and $k_{pp}^{(i)}$ from K_{PRP} ;

Send $\{\alpha_i, k_{pp}^{(i)}\}$ to all the cloud servers;

Receive from servers:

$\{R_i^{(j)} = \sum_{q=1}^r \alpha_i^q * G^{(j)}[\phi_{k_{pp}^{(i)}}(q)] \mid 1 \leq j \leq n\}$

for $(j \leftarrow m + 1, n)$ do

$R^{(j)} \leftarrow R^{(j)} - \sum_{q=1}^r f_{k_j}(s_{I_q, j}) \cdot \alpha_i^q, I_q = \phi_{k_{pp}^{(i)}}(q)$

end for

if $((R_i^{(1)}, \dots, R_i^{(m)}) \cdot P = (R_i^{(m+1)}, \dots, R_i^{(n)}))$ then

Accept and ready for the next challenge.

else

for $(j \leftarrow 1, n)$ do

if $(R_i^{(j)} \neq v_i^{(j)})$ then

return server j is misbehaving.

end if

end for

end if

end procedure

Our token-based approach, while allowing efficient storage correctness validation, does not have this limitation on the number of misbehaving servers. That is, our approach can identify any number of misbehaving servers for $b \leq (m+k)$.

3.3. File Retrieval and Error Recovery

In this paper, we focus on this issue related to retrieval of a file in efficient manner. The tokens of each block which we were generated using precomputational algorithm has been stored in the database. Now we are using homomorphic technique to retrieve entire file or required blocks dynamically. Once user has been sent the requested file to TPA. TPA monitors whether he is authenticated user or not for accessing the file. TPA maintains the file details and tokens (if TPA is not present user will have the details) but not an entire file, TPA requests the file by passing the pre-computed token stored in the database for each block. If this token is same as it is present in cloud server, cloud server will send the requested blocks.

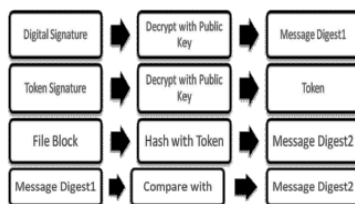
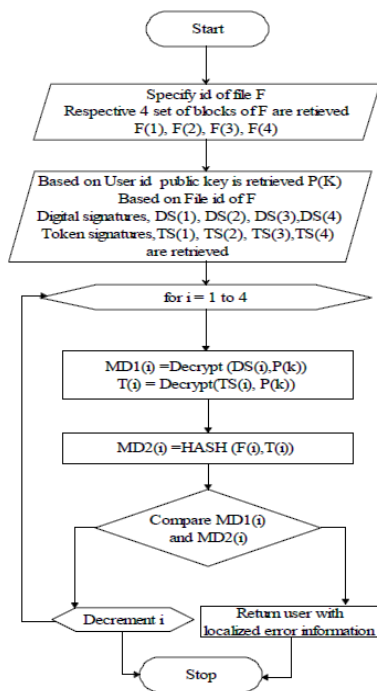


Fig. 2. Operations during File download process [5]

We can easily check whether the file blocks were damaged or not by computing tokens dynamically as follows:

When TPA challenges or requests a block with block indices, cloud server receives this input and it computes the token of that particular block and sends the short signature to TPA. Upon receiving the signature TPA verifies it with the existing token signature. The result of two tokens are same means the block remains same without any effect, otherwise TPA assumes block was modified and it generates a message to cloud server to perform block recovery operation using distributed schemes and erasure coded techniques.

4. Support for Dynamic Data Operation

If we consider 'f' as a file to be stored across cloud servers, some scenarios may arise where users may wish to perform various operations such as updating, deletion and addition at block level. In existing schemes user has to download entire file and then it has to perform operations and again it has to recompute token and digital signatures for entire file which results in overhead for users and is inefficient.

But scheme used in this paper divides file 'f' into blocks, hence user can download only blocks on which it wishes to perform dynamic operation and recompute token and signatures only for those particular blocks not entire file there by providing efficient method for dynamic operations and reducing overhead for users.

4.1. Update operation

In cloud data storage, if user wish to modify any particular block then, it has to specify index of that particular block, then only that particular block is downloaded rather than downloading entire file, only that block is modified and token and signature is recomputed only for that particular modified block and then again uploaded.

4.2. Delete operation

Whenever user wish to delete a particular block then it has to specify the index of that block, download that block and replaces that block with zero or any special character and then recomputed signature for that and uploads back to cloud server.

4.3. Append Operation

Some scenarios may arise where user wish to increase size of data stored across servers by

increasing the number of data blocks. This addition of block will beat the end of the file also known as appending of file blocks, token and signature has to be computed for the newly added block and then uploaded.

5. Related Works

In this section some of the related works are discussed along with their schemes and disadvantages also solution for those disadvantages in the proposed scheme. Some of the related schemes proposed are [6] [7] [8]. Here TPA concept has been proposed. But redundant copy maintenance of file is not included hence does not assure availability of data in case of server failure or corruption of stored cloud data. Since distributed file storage not included, focuses only on single server scenario.

Related works [3] [9] proposed scheme which includes distributed storage that is dividing off files into multiple blocks and storing them randomly across multiple cloud servers, maintaining redundant copy of data to ensure cloud data availability, TPA which checks for cloud data integrity and also error localization performed by TPA to localize at which server the file block has been corrupted. But there is an overhead for users to generate keys. Also does not provide, efficient storage of cloud data, scheme for access of cloud data only for authorized users.

Hence to overcome above specified disadvantages proposed schemes in [5] paper includes efficient storage of cloud data which achieved by storing of data across cloud servers based on their cost and quality and also admin module proposed in [5] paper to achieve authorized access of data stored across cloud servers and reduces the overhead of users of key generation.

But in previous [5] paper in order to perform dynamic operations on data stored across servers entire file has to be downloaded and token and signatures has to be recomputed for entire file, even though the dynamic operations is to be performed at block level, only for particular block which is inefficient one, in order to overcome this scheme proposed in this paper, so that only a particular block can be downloaded on which operations is to be performed. Token and signatures computed only for that particular block not for entire file as in previous scheme there by providing more efficiency.

The work presented in this paper have previously appeared as an extended abstract in [3].

I have revised the paper and add technical details as compared to [3].

The primary improvements are as follows:

- First, we provide the protocol extension for privacy-preserving third-party auditing, and discuss the application scenarios for cloud storage service.
- Second, we add correctness analysis of proposed storage verification design.
- Third, we completely redo all the experiments in our performance evaluation part, which achieves significantly improved result as compared to [3].

6. Conclusion and Future Scope

In cloud computing, IT departments can quickly meet requests for services and time-to-market while mitigating risk and maintaining influence. To the end users: Quick and easy resource sharing, rapid deployment, self-service and the ability to perform chargeback to departments or user groups. Scheme used in this paper, an effective and flexible distributed storage verification scheme with explicit dynamic data support to ensure the correctness and availability of user's data in the cloud. We ensure that the data which was sent to the Cloud Service Provider (CSP) are acknowledged by generating the token dynamically. We will come up with some security algorithms so that the drawbacks are eliminated which in turn can be helpful to increase the security for the end user's data stored on the cloud.

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